

# Homework 3

*Due December 5, 2011*

- Let  $A = \{0, 1\}$  be an alphabet. Define the language  $L$  as follows:
  - $\lambda \in L$ .
  - if  $x \in L$ , then  $0x1 \in L$ .
  - if  $x \in L$  and  $y \in L$ , then  $xy \in L$ .
  - Show that if  $x \in L$ , then  $x$  is a word that has exactly as many 0's as 1's.
  - Show that  $L$  is not regular.
- Prove that the language  $\{x \in \{a, b\}^* \mid n_a(x) \leq n_b(x)\}$  is not regular.
- Let  $A = \{a, b\}$  and let  $L$  be the language  $L = aA^*bA^*b$ . Compute the minimal automaton that will accept  $L$ .
- Let  $G = (\{S, X, Y, Z\}, \{a, b\}, S, P)$  be a context-free grammar whose set of production is

$$P = \{S \rightarrow XaYbZc, S \rightarrow abc, X \rightarrow \lambda, X \rightarrow Ybc, Y \rightarrow \lambda, \\ Y \rightarrow Zac, Z \rightarrow \lambda, Z \rightarrow Xab\}$$

Construct an equivalent grammar without erasure rules.

- Construct a grammar  $G_1$  equivalent to  $G$  given in problem 4 such that  $G_1$  is in Chomsky's normal form.