Heaps and Heap Sorting

- Heaps implemented in an Array
- Heap Sorting
Implementing Heaps in an Array

• Two of the programming steps for a heap using links were complicated
  – Finding the next parent for an add
  – Finding the new value for next after a remove
• Those steps are trivial for a heap in an array
• If the index 0 is the root and “next” is the index for a reference to the next open space:
  – After adding an element, just increment next
  – To remove an element, just decrement next
Implementing Heaps in an Array

• The class header, attributes, and constructor:

```java
public class ArrayHeap<T extends Comparable> {...
{    // not shown as a subclass like in textbook
    private static final int DEFAULT_CAPACITY = 50;
    protected int count;
    protected T[] tree;

    public ArrayHeap()
    {
        count = 0;
        tree = (T[]) new Object[DEFAULT_CAPACITY];
    }
```
Implementing Heaps in an Array

• The add method:

```java
public void addElement(T obj) {
    if (count == tree.length)
        expandCapacity(); // same as for any array

    tree[count++] = obj;

    if (count > 1)
        heapifyAdd();
}
```
Implementing Heaps in an Array

• The heapifyAdd helper method

```java
private void heapifyAdd()
{
    int next = count - 1;
    T temp = tree[next]; // pick up new value

    while ((next != 0) && // move up the tree as needed
           temp.compareTo(tree[(next-1)/2]) < 0 {
        tree[next] = tree[(next-1)/2]; //move parent down
        next = (next - 1)/2;
    }
    tree[next] = temp; // (re)insert new value
}
```
Implementing Heaps in Arrays

• The `removeMin` method:

```java
public T removeMin() {
    // check for empty heap not shown
    T minElement = tree[0];  // start at the root
    tree[0] = tree[count-1];
    heapifyRemove();
    tree[--count] == null;   // kill stale reference

    return minElement;
}
```
Implementing Heaps in Arrays

- The `heapifyRemove` method (My version):

  ```java
  private void heapifyRemove()
  {
      int node = 0;
      T temp = tree[node];
      int next = 0;
      do {
          if (next != 0) { // skip until second+ pass
              tree[node] = tree[next];
              node = next;
          }
      }
  ```
Implementing Heaps in an Array

• The heapifyRemove method (continued):

```java
int left = 2 * node + 1;
int right = 2 * (node + 1);
if (tree[left] == null && tree[right] == null)
    next = count; // force end of loop
else if (tree[right] == null ||
    tree[left].compareTo(tree[right]) < 0)
    next = left;
else
    next = right;
}
while (next < count &&
    tree[next].compareTo(temp) < 0);

tree[node] = temp;
```
Implementing Heaps in Arrays

- I didn’t like the repetition of 8 lines of code in the initialization before the while loop and as 80% of the code in the body of the loop.
- Whenever you see that situation in code, it is a clue that a do-while loop might be better.
- I rewrote the heapifyRemove method as a do-while loop instead of the textbook code that uses a while loop - both code versions perform the same steps in the same order.
Heap Sorting

• If we have a method with a parameter that is an array of type T to be sorted, it can be written to use a heap instead of one of the in-place array sorts we studied in lecture 19.

• We can take each element from the array and put it in a heap.

• Then loop removing the min from the heap and putting each element back into the array in order.
public void heapsort(T[] data) {
    HeapADT<T> temp = new ArrayHeap<T>();
    // copy the array into the heap
    for (T datum : data) { // use a for-each loop
        temp.addElement(datum);
    }
    // place the sorted elements back into the array
    int count = 0;
    while (!temp.isEmpty())
        data[count++] = temp.removeMin();
} // temp goes out of scope
Heap Sorting Performance

• The performance is $2*N*\log N$ or $O(N \log N)$
• That is the same as quicksort or merge sort
• It uses the same amount of extra memory as the merge sort algorithm