

# CS724: Topics in Algorithms

## Problem Set 2

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## Problem 1:

Let  $A \in \mathbb{C}^{m \times n}$  and  $B \in \mathbb{C}^{n \times p}$  be two conformant matrices. Prove that:

- computing the matrix  $G = AB$  using standard matrix multiplication requires  $mnp$  number multiplications;
- if  $C \in \mathbb{C}^{p \times q}$  the computation of the matrix  $D = (AB)C = A(BC)$  by the standard method, the first modality  $D = (AB)C$  requires  $mp(n + q)$  multiplications, while the second,  $D = A(BC)$  requires  $nq(m + p)$  multiplications.



## Solution 1:

Each element  $g_{ij}$  of  $G \in \mathbb{C}^{mp}$  requires  $n$  multiplications because

$$g_{ij} = \sum_{k=1}^n a_{ik} \cdot b_{kj}.$$

Since there are  $mp$  such elements, the total number of multiplications is  $mpn$ .

Note that  $AB \in \mathbb{R}^{m \times p}$  and  $BC \in \mathbb{R}^{n \times q}$ . Therefore,  $(AB)C$  requires  $mp(n + q)$  multiplications, and  $A(BC)$  requires  $nq(m + p)$  multiplications. A judicious organization of computations would compare these numbers. If  $mp(n + q) < nq(m + p)$ , then  $(AB)C$  is preferable to  $A(BC)$  because it would involve fewer multiplications.



## Problem 2:

Let  $A, B$  be two matrices in  $\mathbb{C}^{n \times n}$ . Suppose that  $B = C + D$ , where  $C$  is a Hermitian matrix and  $D$  is a skew-Hermitian matrix. Prove that if  $A$  is Hermitian and  $AB = BA$ , then  $AC = CA$  and  $AD = DA$ .



## Solution 2:

Since  $C$  is Hermitian we have  $C^H = C$ . Since  $D$  is skew-Hermitian,  $D^H = -D$ . Let  $A$  be such that  $A^H = A$  and  $AB = BA$ .

$AB = BA$  is equivalent to  $A(C + D) = (C + D)A$ .

Also,  $AB = BA$  implies  $B^H A = AB^H$  because  $A$  is Hermitian. Since  $B^H = C^H + D^H = C - D$ , this amounts to  $A(C - D) = (C - D)A$ . Since we also have  $A(C + D) = (C + D)A$  by adding the last two equalities we obtain  $2AC = 2CA$  and, by subtracting then, we have  $2AD = 2DA$ , which yield the conclusion.



## Problem 3:

Recall that  $J_{n,n} \in \mathbb{R}^{n \times n}$  is the complete  $n \times n$  matrix, that is the matrix having all components equal to 1. Prove that for every number  $m \in \mathbb{N}$  and  $m \geq 1$  we have

$$J_{n,n}^m = n^{m-1} J_{n,n}.$$



## Solution 3:

Observe that  $J_{n,n}^2 = nJ_{n,n}$ .

The proof is by induction on  $m$ . The base step,  $m = 1$  is immediate.

Suppose this holds for  $m$ . Then

$$\begin{aligned} J_{n,n}^{m+1} &= J_{n,n} J_{n,n}^m = J_{n,n} \cdot n^{m-1} J_{n,n} \\ &= \text{(by inductive hypothesis)} \\ &= n^m J_{n,n}. \end{aligned}$$



## Problem 4:

Let  $A \in \mathbb{R}^{n \times n}$  and  $B \in \mathbb{R}^{n \times n}$  be two matrices. Prove that if  $AB$  is invertible, then both  $A$  and  $B$  are invertible.





## Solution 4:

Since  $AB$  is invertible we have  $AB(AB)^{-1} = I_n$ . Thus,  $A$  is invertible and  $A^{-1} = B(AB)^{-1}$ . Similarly, since  $(AB)^{-1}AB = I_n$  it follows that  $B$  is invertible and  $B^{-1} = (AB)^{-1}A$ .



## Problem 5:

Let  $A = (a_{ij})$  be an  $(m \times n)$ -matrix of real numbers. Prove that

$$\max_j \min_i a_{ij} \leq \min_i \max_j a_{ij}$$

(the *minimax inequality*).



## Solution 5:

Observe that  $a_{ij_0} \leq \max_j a_{ij}$  for every  $i$  and  $j_0$ , so  $\min_i a_{ij_0} \leq \min_i \max_j a_{ij}$ , again for every  $j_0$ . Thus,  $\max_j \min_i a_{ij} \leq \min_i \max_j a_{ij}$ .



## Problem 6:

Let  $\mathbf{x}, \mathbf{y} \in \mathbb{R}^n$  be two vectors such that  $\mathbf{x}'\mathbf{y} \neq 0$  and  $\mathbf{xy}' \neq O_{n,n}$ .

Prove that both matrices  $\mathbf{x}'\mathbf{y}$  and  $\mathbf{xy}'$  have rank 1.

Note that if  $\mathbf{x} \in \mathbb{R}^m$  (or  $\mathbb{R}^{m \times 1}$ ) and  $\mathbf{y} \in \mathbb{R}^n$  (or  $\mathbb{R}^{n \times 1}$ ) and  $m \neq n$  the multiplication  $\mathbf{x}'\mathbf{y}$  can not be performed, but  $\mathbf{xy}'$  is feasible and the result is the same (matrix  $\mathbf{xy}'$  has rank 1).



## Solution 6:

Note that  $\mathbf{x}'\mathbf{y}$  is a number:

$$\mathbf{x}'\mathbf{y} = x_1y_1 + \cdots + x_ny_n,$$

and therefore has rank 1.

On other hand,

$$\mathbf{xy}' = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix} \mathbf{y}' = \begin{pmatrix} x_1\mathbf{y}' \\ \vdots \\ x_n\mathbf{y}' \end{pmatrix},$$

which implies that  $\mathbf{xy}'$  also has rank 1 because every row of this matrix is a multiple of  $\mathbf{y}'$ .

