

Course Info

Lecture Hours

Mon and Wed , 8:30-9:45pm

Office Hours

- Mon & Wed 9:00-10:30am
- By appointment (send email)

Class URL

- http://www.cs.umb.edu/~gghinita/cs430/
- http://www.cs.umb.edu/~gghinita/cs630/

3

Textbook & Recommended Readings

Textbook

 Database Management Systems, 3rd Edition by Ramakrishnan and Gehrke



Other recommended texts

- Database System Concepts, Silberschatz, Korth and Sudarshan, 6th Edition
- Database Principles, Programming, and Performance, P. E. O'Neil and E. J. O'Neil
- > Other resources will be posted in the links section of the site

Prerequisites

- Data Structures and Algorithms
 - CS310
- Programming
 - CS240
- Discrete Math
- Familiarity with UNIX OS
 - Exercises will be executed on Oracle 12G server running on a Unix machine in the CS dept (DBS3 Machine)

▶ 5

Grading

• 4

- Final exam (40%) open book
- Midterm (30%) open book (Wed March 23th)
- 6 homework assignments
 - 5% each
- Assignments for CS630 will have additional questions
- Assignments are individual submit your own work only!
- No plagiarism! See student code of conduct
- Lecture attendance is mandatory

▶ 6

Course Materials

Class URL

- http://www.cs.umb.edu/~gghinita/cs430/
- http://www.cs.umb.edu/~gghinita/cs630/
- Blackboard
 - Discussion forums
- Make sure you create Unix course accounts, and that you enroll these accounts for 630 ("apply" procedure)
- ▶ 7

University Policies

Student Conduct: Students are required to adhere to the University Policy on Academic Standards and Cheating, to the University Statement on Plagiarism and the Documentation of Written Work, and to the Code of Student Conduct as delineated in the University Catalog and Student Handbook. The Code is available online at:

https://www.umb.edu/life_on_campus/dean_of_students/student_conduct

Accommodations: Section 504 of the Americans with Disabilities Act of 1990 offers guidelines for curriculum modifications and adaptations for students with documented disabilities. If applicable, students may obtain adaptation recommendations from the Ross Center for Disability Services, CC-UL Room 211, (617-287-7430). The student must present these recommendations and discuss them with each professor within a reasonable period, preferably by the end of Drop/Add period.

Course Overview

- Relational Data Model
- Relational Algebra
- Structured Query Language
- > The most important part of the course
- Conceptual design the ER model
- Database application development
- Java, PL/SQL
- Design Theory
- Database Security

▶ 9

What is a DBMS? • Specialized software that provides:

- Uniform and transparent access to data
 - Application-independence
 - Application/user is oblivious to internal data organization
 - Data organization may change, but applications need not change
- Efficient access to data
 - Fast search capabilities, indexing
- Data consistency
 - E.g., cannot delete student record if grade records still in DBMS
- Concurrent access to data

A bit of history ...

- Persistent storage and recovery from failure
- Security

▶ 10

▶ 8

Why study databases?

> Databases are ubiquitous

- Behind all web service providers there is a DBMS
 - Most often a very large-scale one
- Corporations use DBMS for business processes, HR, etc
- Scientific computing relies on very large amounts of data
 - Humane genome data
 - Biochemistry data (protein sequences)
 - Astronomy data
 - High-energy physics
- DBAs are very well-paid!
 - And even in other IT areas, DBMS skills are a must

► 11

- First data stores were file systems
 Does not conform to transparency and uniformity desiderata
 Search (within file) most often linear
 Not portable
 Doesn't handle concurrency properly
 Sequential access only
 Early DBMS appeared in the 60s
 Driven by banking and airline industry
 Relatively small record size, and many concurrent accesses
 - Two prominent models: hierarchical model (tree) and network model (graph)
 - Lack of support for high-level query languages

12

A bit of history (contd.)

Relational Databases

- Major breakthrough, paper written by Codd (1970)
- Relations (tables) with rows (records) and columns (fields)
 Relationships and constraints among tables
- Structured Query Language (SQL): high-level, declarative
 Data definition/ manipulation language
- Fast search use of index structures
- > Data access language independent from internal organization

Newer paradigms

- Object-oriented and multimedia DB
- Data Stream Management Systems (DSMS)
- MapReduce

13





Data Model Structure of data Relational model uses tables Programming languages deal with arrays, collections, etc Operations on the data Queries: operations that retrieve information Modifications: operations that change data Constraints

- Domain constraints (the simplest): e.g., age must be numeric
- Other constraints: each student has unique matriculation #
- Prominent Data Models
- Relational model
- > Object-relational model, semi-structured model (XML), E-R

Relational Model > Relational database: a set of relations > Relation: > two-dimensional table, with rows and columns #Rows = cardinality #Columns= degree (or arity) > Each row represents an entity > A student, a course, a movie > Each column represents a property of the entity > Student age, student matriculation #, student gpa > Column values are atomic (e.g., integer or string) within given domain > Rows are also called tuples or records; columns are also called fields or attributes

53666	Iones	iones@cs	18	34
53688	Smith	smith@eecs	18	3.2
53650	Smith	smith@math	19	3.8
Car	dinality	= 3, Degree =	5	

▶ 17

▶ 18



More about Relations

- Relations are sets of tuples
 Sets are NOT ordered
 - Do NOT retrieve by order number, but by content!

_ . . .

- Relation Instance
 Contents of a relation may change over time
 - Tuples are added/deleted/modified
 - E.g., Students join or leave the university
 - Instance represents set of tuples at a certain point in time

Schemas may change too

- > Although this is not frequent in practice
- Changing schema is very expensive

> 20

53666 Jones jones@cs 18		ine rogin	na	S1d
	18	es jones@cs	Joi	53666
53688 Smith smith@eecs 18	18	th smith@eecs	Sn	53688
53650 Smith smith@math 19	19	th smith@math	Sn	53650

n	other	Instan	ce of "Stude	ents"	
	sid	name	login	age	gpa
	53666	Jones	jones@cs	18	3.4
	53688	Smith	smith@eecs	18	3.2
	53650	Smith	smith@math	19	3.8
	53660	Korth	korth@math	22	3.6
	-		•		
	C	ardinali	ity = 4, Degree	= 5	

Keys A key of a relation is a set of fields K such that: No two distinct tuples in ANY relation instance have same values in all key fields, and No subset of K is a key (otherwise K is a superkey) Key may not be unique Multiple candidate keys may exist One of the keys is chosen (by DBA) to be the primary key Keys are shown <u>underlined</u> in schema In the relational model, duplicate tuples do not exist! But most DBMS implementations do allow duplicates

- Keys constraints must be set by DBA to avoid duplicates
- 23

Example of Keys

Students(sid:string, name: string, login: string, age: integer, gpa: real)

sid is a key; {sid, name} is a superkey

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@eecs	18	3.2
53650	Smith	smith@math	19	3.8

- In practice, it is not easy to know when there exists a unique attribute combination in the data (e.g., names)
- > artificial keys are created: student ID, customer ID, etc.
- SSN is also often used for keys

24

22