

Introduction to Programming in Python

Algorithms and Data Structures: Analysis of Algorithms

Outline

① Performance

② Time Complexity

③ Space Complexity

Performance

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For example, if $f(n) = 31n^2 + 78n + 42$, then $T(n) = n^2$

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```
>_ ~/workspace/dsa/programs
```

```
$ cat ../data/1Kints.txt
324110
-442472
...
745942
$ /usr/bin/time --format='%e seconds' python3 triplesum.py ../data/1Kints.txt
70
0.7 seconds
$ /usr/bin/time --format='%e seconds' python3 triplesum.py ../data/2Kints.txt
528
5.9 seconds
```

Time Complexity

Time Complexity

</> triplesum.py

```
from instream import InStream
import stdio
import sys

def main():
    inStream = InStream(sys.argv[1])
    a = inStream.readAllInts()
    stdio.writeln(count(a))

def count(a):
    n = len(a)
    count = 0
    for i in range(0, n):
        for j in range(i + 1, n):
            for k in range(j + 1, n):
                if a[i] + a[j] + a[k] == 0:
                    count += 1

    return count

if __name__ == '__main__':
    main()
```


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1K	0.28s
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4K	14.06s
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[B]
[C]
[D]
[E]

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Statement Block	Time	Frequency	Total Time
[A]	t_4	1	t_4
[B]	t_3	n	$t_3 n$
[C]	t_2	$\binom{n}{2}^1 = n^2/2 - n/2$	$t_2(n^2/2 - n/2)$
[D]	t_1	$\binom{n}{3} = n^3/6 - n^2/2 + n/3$	$t_1(n^3/6 - n^2/2 + n/3)$
[E]	t_0	\times (depends on input)	$t_0 \times$

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Grand total: $f(n) = (t_1/6)n^3 + (t_2/2 - t_1/2)n^2 + (t_1/3 - t_2/2 + t_3)n + t_4 + t_0 \times$

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Grand total: $f(n) = (t_1/6)n^3 + (t_2/2 - t_1/2)n^2 + (t_1/3 - t_2/2 + t_3)n + t_4 + t_0 \times$

Running time: $T(n) = n^3$

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Running time classifications

Name	$T(n)$	Code Description	Example
constant	1	statement	increment the i th element in an array
logarithmic	$\log n$	divide and discard	binary search
linear	n	loop	find the maximum
linearithmic	$n \log n$	divide and conquer	merge sort
quadratic	n^2	double loop	check all ordered pairs
cubic	n^3	triple loop	check all ordered triples
exponential	2^n	exhaustive search	check all subsets

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Sizes of built-in objects on a typical system

Object	Size in Bytes
integer	24
float	24
boolean	24
string of n characters	$40 + n$
list of n integers	$72 + 8n + 24n = 72 + 32n$
m -by- n list of integers	$72 + 8m + m(72 + 32n) = 72 + 80m + 32mn$
user-defined	hundreds of bytes, at least