## 1 Exercises

Exercise 1. Consider the *j*-- program sum shown below:

```
package pass;
import java.lang.Integer;
import java.lang.System;
public class Sum {
    private static String MSG = "SUM = ";
    private int n;
    public Sum(int n) {
        this.n = n:
    }
    public int compute() {
        int sum = 0, i = n;
        while (i > 0) {
            sum += i--:
        }
        return sum;
    }
    public static void main(String[] args) {
        int n = Integer.parseInt(args[0]);
        Sum sum = new Sum(n):
        System.out.println(MSG + sum.compute());
    7
}
```

How does JVM bytecode generation (JCompilationUnit.codegen()) for the program work?

Exercise 2. Suppose 1hs and rhs are boolean expressions. How does *j*-- generate code for the following statements?

```
a.
1 boolean x = lhs && rhs;
b.
1 if (lhs && rhs) {
2 then_statement
3 } else {
4 else_statement
5 }
```

c.

```
1 | while (lhs && rhs) {
2 statement
3 }
```

**Exercise 3.** Suppose  ${\tt x}$  is an object and  ${\tt y}$  is an integer field within.

a. What is the JVM bytecode generated for the following statement? How does the runtime stack evolve as the instructions are executed?

1 ++x.y;

b. If z is also an integer, what is the JVM bytecode generated for the following statement? How does the runtime stack evolve as the instructions are executed?

1 z = ++x.y;

Exercise 4. How is code generated for the expression "The first perfect number is " + 6?

Exercise 5. How is code generated for casts?

Exercise 6. How would you generate JVM bytecode for the do-while statement, ie, implement codegen() in JDoWhileStatement.java?

## 2 Solutions to Exercises

Solution 1. Consult sections 5.2 – 5.6 of our text.

## Solution 2.

```
a.
           lhs code
           branch to Target on false
           rhs code
           branch to Target on false
           push 1 on stack
           goto End
   Target: push 0 on stack
   End:
           . . .
b.
           lhs code
           branch to Target on false
           rhs code
           branch to Target on false
           then_statement code
           goto End
   Target: else statement code
   End:
           . . .
C. Test:
           lhs code
           branch to Target on false
           rhs code
           branch to Target on false
           body code
           goto Test
   Target:
```

Solution 3. We use table on slide 26 from the JVM Code Generation chapter.

```
\mathbf{a}. Bytecode:
```

aload x' dup . getfield y iconst\_1 iadd putfield y Runtime stack (right to left is top to bottom): x | x | x | x | y | x | y | 1 | x | y+1 b. Bytecode: aload x dup getfield y iconst\_1 iadd dup\_x1 putfield y Runtime stack (right to left is top to bottom): | x | | x | x | x | y | x | y | 1

| x | y+1 | y+1 | x | y+1 | y+1

Solution 4. Since the left-hand-side expression of + is a string, the operation denotes string concatenation, and is represented in the AST as a <code>JstringConcatenationOp</code> object. The <code>codegen()</code> method therein does the following:

- 1. Creates an empty string buffer, ie, a stringBuffer object, and initializes it.
- 2. Appends the string "The first perfect number is " to the buffer using stringBuffer's append(String x) method.
- 3. Appends the integer value 6 to the buffer using stringBuffer's append(int x) method.

4. Invokes the tostring() method on the buffer to produce a string on the runtime stack.

Solution 5. Analysis determines both the validity of a cast and the necessary converter, which encapsulates the code generated for the particular cast. Each Converter implements a method codegen(), which generates any code necessary to the cast. Code is first generated for the expression being cast, and then for the cast, using the appropriate converter.

## Solution 6.

2

3

4

5

6

```
public void codegen(CLEmitter output) {
   String bodyStart = output.createLabel();
   output.addLabel(bodyStar);
   body.codegen(output);
   condition.codegen(output, bodyStart, true);
}
```