CS450 High Level Languages  (section 2)

Abstraction

Wednesday, October 11, 2023
Logistics

• HW 3 out
  • due: Sun 10/15 11:59 pm EST

• HW 2 (and other) grades returned
  • Use GradeScope re-grade request for complaints or questions
CS 450 so far ...

This class teaches:

• a high-level programming “process”
• i.e., a design recipe for creating clean, readable programs

• How to do well: learn and follow the process

• How to not do well: focus only on “getting the code working”
“Perhaps you thought that “getting it working” was the first order of business for a professional developer.

I hope by now, however, that this book has disabused you of that idea.

The functionality that you create today has a good chance of changing in the next release, but the **readability of your code** will have a profound effect on all the changes that will ever be made.”

— Robert C. Martin,  
**Clean Code: A Handbook of Agile Software Craftsmanship**
HW 2 recap

Many submissions only focused on: “getting the code working”

Many submissions ignored:
- all other steps of programming design recipe
- style guide
- Other instructions in hw

This hw will be graded accordingly:
- correctness (9 pts)
- design recipe (20 pts)
  - style (5 pts)
  - README (1 pt)
Total: 35 points
HW 2 solution (part 2)

“world state” now needs two “posn”s!

- To practice reading code, save the big-bang program from lecture 6 to a file named hw2-bigbang.rkt, and change it so when the user clicks the left-mouse button, a copy of the image is pinned to the canvas. (The original image should continue to move with the cursor.) Subsequent clicks should move the pinned image so that only one image is pinned at any time.
WorldState is a:
(struct world [cursor pinned])
where
cursor : Posn - IMG location that moves with mouse cursor
pinned : MaybePosn - IMG location that is pinned (if there is one)

Posn is a
(struct posn [x y])
where
x: Integer - represents x coordinate in big-bang animation
y: Integer - represents y coordinate in big-bang animation

MaybePosn is a
- Posn
- NO-POSN

constants and predicates for MaybePosn
(define NO-POSN #false)
(define (no-posn? x) (equal? x NO-POSN))
(define (maybe-posn? x) (or (posn? x) (no-posn? x)))
Mouse handler

;; mouse-handler : WorldState Coordinate Coordinate MouseEvent -> WorldState
;; Returns a WorldState where:
;; - "cursor" = current mouse loc
;; - "pinned" = current mouse loc (if mevt = "button-down")

;; TEMPLATE fn for MouseEvent : MouseEvent -> ???
(define (mevt-fn mevt)
  (cond
    [(string=? mevt "button-down") ....]
    .... ])

Function splitting rule:
One (data definition processing) task, one function

When a function has more than one argument, you choose which template to use
Mouse handler

;;; mouse-handler : WorldState Coordinate Coordinate MouseEvent -> WorldState
;;; Returns a WorldState where:
;;; - "cursor" = current mouse loc
;;; - "pinned" = current mouse loc (if mevt = "button-down")

(define (mouse-handler w x y mevt)
  (cond
    [(string=? mevt "button-down")
      (world (posn x y) (posn x y))]
    [else
      (world (posn x y) (world-pinned w))])
)

(define current-mouse-pos (posn x y))

(define (mouse-handler w x y mevt)
  (cond
    [(string=? mevt "button-down")
      (world current-mouse-pos current-mouse-pos)]
    [else
      (world current-mouse-pos (world-pinned w))])))

This field always the same
Mouse handler

;; mouse-handler : WorldState MouseEvent -> WorldState
;; Returns a WorldState where:
;; - "cursor" = current mouse loc
;; - "pinned" = current mouse loc (if mevt = "button-down")
(define (mouse-handler w x y mevt)
    (define current-mouse-pos (posn x y))
    (define (mouse-handler w x y mevt)
    (define current-mouse-pos (posn x y))
    (world
        current-mouse-pos
        (cond
            [(string=? mevt "button-down") current-mouse-pos]
            [else (world-pinned w)])))

WAIT! Function split rule is:
One (data definition processing) task, one function

Do we need a separate (world-processing) function?
On when to create a new function ...

“The first rule of functions is that they should be small.

The second rule of functions is that they should be smaller than that.”

— Robert C. Martin,
Clean Code: A Handbook of Agile Software Craftsmanship

In this class:
create one function per (data definition processing) task
Mouse handler

(define (mouse-handler w x y mevt)
  (define current-mouse-pos (posn x y))
  (world
    current-mouse-pos
    (cond
      [(string=? mevt "button-down") current-mouse-pos]
      [else (world-pinned w)])))

WAIT! Function split rule is:
One (data definition processing) task, one function

Do we need a separate (world-processing) function?

YES. (But in this case, it would just get the world-pinned value)
Render world

;;; render-world : WorldState -> Image
;;; draws IMG at the "cursor" and "pinned" posn (if there is one)

;;; TEMPLATE fn for WorldState: WorldState -> ???
(define (world-fn w)
  .... (world-cursor w) ..... (world-pinned w) ....)

These are posns. Should we also extract posn fields here?

NO. They should be handled by a “posn” function (which follows “posn” template)

create one function per (data definition processing) task
Render world

;; render-world : WorldState -> Image
;; draws IMG at the "cursor" and "pinned" posn (if there is one)

;;; TEMPLATE fn for WorldState: WorldState -> ???
(define (world-fn w)
  .... (world-cursor w) .... (world-pinned w) ....)

(define (render-world w)
  render-posn
  (world-pinned w)
  (render-posn
    (world-cursor w)
    EMPTY-SCENE)))

Wish list:
render-posn
maybe-render-posn

But this is a “maybe” posn
Render posn

;;; render-posn : Posn Image -> Image
;;; draws IMG into given image at given posn;

;;; TEMPLATE fn for Posn fn: Posn-> ???
(define (posn-fn p)
    .... (posn-x p) .... (posn-y p) ....)

(define (render-posn p img)
    (place-image IMG (posn-x p) (posn-y p) img))
Maybe Render posn

`; maybe-render-posn : MaybePosn Image -> Image
`; draws IMG into given image at given posn;
`; if posn is NO-POSN, return img unchanged
`

`; TEMPLATE fn for MaybePosn fn: MaybePosn-> ??
(define (maybeposn-fn p)
  (cond
   [(posn? p) .... (posn-x p) .... (posn-y p) ....)]
   [(no-posn? p) ....]])

(define (maybe-render-posn p img)
  (cond
   [(posn? p) (place-image IMG (posn-x p) (posn-y p) img)]
   [else img]]))

Same as render-posn
Maybe Render posn

;; maybe-render-posn : MaybePosn Image -> Image
;; draws IMG into given image at given posn;
;; if posn is NO-POSN, return img unchanged

;; TEMPLATE fn for MaybePosn fn: MaybePosn-> ???
(define (maybe-posn-fn p)
  (cond
   [(posn? p) .... (posn-x p) .... (posn-y p) ....]
   [(no-posn? p) ....]))

(define (maybe-render-posn p img)
  (cond
    [(posn? p) (render-posn p img)]
    [else img])))
Render world

;; render-world : WorldState -> Image
;; draws IMG at the "cursor" and "pinned" posn (if there is one)

(define (render-world w)
  (maybe-render-posn
   (world-pinned w)
   (render-posn
    (world-cursor w)
    EMPTY-SCENE)))

Wish list:
render-posn
maybe-render-posn
Render world: alternate choice

;;; render-world : WorldState -> Image
;;; draws IMG at the "cursor" and "pinned" posn (if there is one)

```
(define (render-world w)
  (maybe-render-posn
   (world-pinned w)
   (maybe-render-posn
    (world-cursor w)
    EMPTY-SCENE)))
```

Wish list:
- render-posn
- maybe-render-posn

Works for both!

Using the same function could be more readable when there are an arbitrary number of balls...
HW 2 recap - other points

- GitHub repos must be added to cs450f23 organization (not your own account)
  - Otherwise I can't see it
  - (see directions in hw)

- Git commit messages must be meaningful

Do not...
- Add to temp files to repo
- Leave commented out code
- Leave “TODO”s in code
Lists and List Functions Review
Previously

Racket List Data Definition Example

```scheme
;; A ListofInt is one of
;; - empty
;; - (cons Int ListofInt)
```

- **Empty (base) case**
- **Non-empty (recursive) case**

**cons = “node”**

**Recursive!**
(Using a definition to define itself)

**TEMPLATE??**

(how can we use a list of ints to define a list of ints?!!?)

**Recursion** is a valid concept (from math), but only if there is both
- A **base** case
- A **recursive** case
Previously

Racket Recursive List Fn Template

;; A ListofInt is one of
;; - empty
;; - (cons Int ListofInt)

;; TEMPLATE for list-fn
;; list-fn : ListofInt -> ???
(define (list-fn lst)
  (cond
    [(empty? lst) ....]
    [(cons? lst) .... (first lst) ....
    .... (list-fn (rest lst)) ....]))
Racket Recursive List Fn: \texttt{inc-list}

\begin{verbatim}
;; TEMPLATE for list-fn
;; list-fn : ListofInt -> ???
(define (list-fn lst)
  (cond
   [(empty? lst) ....]
   [(cons? lst) .... (first lst) .... .... (list-fn (rest lst)) ....]))
\end{verbatim}
inc-list : ListofInt -> ListofInt
;; increments each list element by 1
(define (inc-lst lst)
  (cond
    [(empty? lst) ....]
    [(cons? lst) .... (first lst) ....
     .... (inc-lst (rest lst)) ....])))
Racket Recursive List Fn: inc-list

;; inc-list : ListofInt -> ListofInt
;; increments each list element by 1
(define (inc-list lst)
  (cond
    [(empty? lst) empty]
    [(cons? lst) .... (first lst) .... .... (inc-list (rest lst)) ....])))
Racket Recursive List Fn: inc-list

;;; inc-list : ListofInt -> ListofInt
;;; increments each list element by 1
(define (inc-list lst)
  (cond
   [(empty? lst) empty]
   [else .... (add1 (first lst)) .... .... (inc-list (rest lst)) ....])))
Racket Recursive List Fn: **inc-list**

```scheme
;; inc-list : ListofInt -> ListofInt
;; increments each list element by 1
(define (inc-lst lst)
  (cond
   [(empty? lst) empty]
   [else (cons (add1 (first lst))
               (inc-lst (rest lst)))]))
```
Multi-ball Animation

Design a **big-bang** animation that:

- **Start**: a single ball, moving with random x and y velocity
- **On a click**: add a ball at random location, with random velocity

;; A WorldState is ... a *list* of balls!
Ball

(struct ball [x y xvel yvel] #:transparent)

where

x: XCoord - represents x coordinate of ball center in animation
y: YCoord - represents y coordinate of ball center in animation
xvel: Integer - represents x velocity, where
    positive = to the right, negative = to the left
yvel: Integer - represents y vel, where
    positive = down, negative = up

A ListofBall is one of
- empty
- (cons Ball ListofBall)

A WorldState is a ListofBall
next-world

;;; next-world : WorldState -> WorldState
;;; Updates position of all balls by one tick
(define (next-world w)
  (cond
    [(empty? w) ....]
    [else .... (first w) ....
     .... (next-world (rest w)) ....]])
next-world

;;; next-world : WorldState -> WorldState
;;; Updates position of all balls by one tick
(define (next-world w)
  (cond
   [(empty? w) empty]
   [else .... (first w) ....
    .... (next-world (rest w)) ....]]))

(check-equal? (next-world (list (make-ball 0 0 1 1)))
  (list (next-ball (make-ball 0 0 1 1))))
next-world

;; next-world : WorldState -> WorldState
;; Updates position of all balls by one tick
(define (next-world w)
  (cond
   [(empty? w) empty]
   [else .... (next-ball (first w)) ....
      .... (next-world (rest w)) ....]])
next-world

```scheme
;; next-world : WorldState -> WorldState
;; Updates position of all balls by one tick
(define (next-world w)
  (cond
   [(empty? w) empty]
   [else (cons (next-ball (first w))
               (next-world (rest w)))]))
```
next-world

```scheme
;;; next-world : ListofBall -> ListofBall
;;; Updates position of all balls by one tick
(define (next-world lst)
  (cond
   [(empty? lst) empty]
   [else (cons (next-ball (first lst))
                (next-world (rest lst)))]))```
Comparison

;; inc-lst: ListofInt -> ListofInt
;; Returns list with each element incremented
(define (inc-lst lst)
  (cond
   [(empty? lst) empty]
   [else (cons (add1 (first lst))
                (inc-lst (rest lst)))]))

;; next-world : ListofBall -> ListofBall
;; Updates position of each ball by one tick
(define (next-world lst)
  (cond
   [(empty? lst) empty]
   [else (cons (next-ball (first lst))
                (next-world (rest lst)))]))
Abstraction: Common List Function #1

;; lst-fn1: (?? -> ??) Listof?? -> Listof??
;; Applies the given fn to each element of given lst
(define (lst-fn1 fn lst)
  (cond
   [(empty? lst) empty]
   [else (cons (fn (first lst))
               (lst-fn1 (rest lst)))]))

(define (inc-lst lst) (lst-fn1 add1 lst)
(define (next-world lst) (lst-fn1 next-ball lst)
Abstraction: Common List Function #1

;; lst-fn1: (X -> X) ListofX -> ListofX
;; Applies the given fn to each element of given lst

(define (lst-fn1 fn lst)
  (cond
   [(empty? lst) empty]
   [else (cons (fn (first lst))
              (lst-fn1 (rest lst)))]))

(define (inc-lst lst) (lst-fn1 add1 lst)
(define (next-world lst) (lst-fn1 next-ball lst))
Abstraction: Common List Function #1

;;; lst-fn1: (X -> Y) ListofX -> ListofY
;;; Applies the given fn to each element of given lst

(define (lst-fn1 fn lst)
  (cond
   [(empty? lst) empty]
   [else (cons (fn (first lst))
                (lst-fn1 (rest lst)))]))

(define (inc-lst lst) (lst-fn1 add1 lst)
(define (next-world lst) (lst-fn1 next-ball lst))
Abstraction Data Definitions

;; A ListofInt is one of
;; - empty
;; - (cons Int ListofInt)

;; A ListofBall is one of
;; - empty
;; - (cons Ball ListofBall)

;; A Listof<X> is one of
;; - empty
;; - (cons X Listof<X>)

To use this abstract data definition, must instantiate \( X \) with a concrete data definition

(concrete = opposite of abstract)

NOTE: this shows why our Compound data predicates should be “shallow” checks, i.e., list?

Makes abstraction easier
Abstract Data Defs common in every PL

```cpp
#include<iostream>
#include<vector>
using namespace std;

int main()
{
    vector<int> v;

    for (int i = 1; i <= 10; i++)
    {
        v.push_back(i);
    }

    cout << "Size : " << v.size();

    v.resize(7);

    cout << "\nAfter resizing it becomes : " << v.size();
}```
A `Posn` is a

```
(struct posn [x y])
```

where

```
;; x: Integer - represents x coordinate in big-bang animation
;; y: Integer - represents y coordinate in big-bang animation
```
Common List Function #1

;;; lst-fn1: (X -> Y) Listof<X> -> Listof<Y>
;;; Applies the given fn to each element of given lst

(define (lst-fn1 fn lst)
  (cond
    [(empty? lst) empty]
    [else (cons (fn (first lst))
                (lst-fn1 (rest lst)))]))

(define (inc-lst lst) (lst-fn1 add1 lst)
(define (next-world lst) (lst-fn1 next-ball lst))
Common List Function #1: \texttt{map}

\begin{verbatim}
;;; map: (X -> Y) Listof<X> -> Listof<Y>  
;;; Applies the given fn to each element of given lst

(define (map fn lst)
  (cond
    [(empty? lst) empty]
    [else (cons (fn (first lst))
                 (map (rest lst))))]))

(define (inc-lst lst) (map add1 lst)
(define (next-world lst) (map next-ball lst))
\end{verbatim}
Common List Function #1: `map`

```
(map proc lst ...) → list?
proc : procedure?
lst : list?
```

`map`: (A B C ... -> Z) Listof<A> Listof<B> Listof<C> ... -> Listof<Z>

;; Applies the given fn to elements (at same index) of given lsts

```
(check-equal? (map + (list 1 2 3) (list 4 5 6))
  (list 5 7 9))
```
Common List Function #2: ???
Previously

Racket Recursive List Fn Example: sum

;; TEMPLATE for list-fn
;; list-fn : ListofInt -> ???
(define (list-fn lst)
  (cond
   [(empty? lst) ....]
   [(cons? lst) .... (first lst) ....
   .... (list-fn (rest lst)) ....])))
Previously

Racket Recursive List Fn Example: sum

;; Returns sum of list of ints
;; sum-lst: ListofInt -> Int
(define (sum-lst lst)
  (cond
    [(empty? lst) 0]
    [else (+ (first lst)
              (sum-lst (rest lst)))]))
Render World: ListofBall edition

;; render-world : ListofBall -> Image
;; Draws the given world as an image by overlaying each ball, ;; at its position, into an initially empty scene

(define (render-world lst)
  (cond
   [(empty? lst) ....]
   [else .... (first lst) .... (render-world (rest lst)) ....])))
Render World: ListofBall edition

;;; render-world : ListofBall -> Image
;;; Draws the given world as an image by overlaying each ball, ;; at its position, into an initially empty scene

(define (render-world lst)
  (cond
   [(empty? lst) EMPTY-SCENE]
   [else .... (first lst).... (render-world (rest lst)) ....])))
Render World: ListofBall edition

;;; render-world : ListofBall -> Image
;;; Draws the given world as an image by overlaying each ball, at its position, into an initially empty scene

(define (render-world lst)
  (cond
    [(empty? lst) EMPTY-SCENE]
    [else (place-ball (first lst) (render-world (rest lst)))]))

;;; place-ball : Ball Image -> Image
;;; Draws a ball, using its pos as the offset, into the given image
(define (place-ball b scene)
  (place-image BALLIMG (ball-x b) (ball-y b) scene))
Comparison #2

;;; sum-lst: ListofInt -> Int
(define (sum-lst lst)
  (cond
   [(empty? lst) 0]
   [else (+ (first lst)
             (sum-lst (rest lst)))]))

;;; render-world: ListofBall -> Image
(define (render-world lst)
  (cond
   [(empty? lst) EMPTY-SCENE]
   [else (place-ball (first lst)
                   (render-world (rest lst)))]))
;;; list-fn2 : (X Y -> Y) Y Listof<X> -> Y

(define (list-fn2 fn initial lst)
  (cond
   [(empty? lst) initial]
   [else (fn (first lst) (list-fn2 fn initial (rest lst)))]))

;;; sum-lst: ListofInt -> Int
(define (sum-lst lst) (list-fn2 + 0 lst))

;;; render-world: ListofBall -> Image
(define (render-world lst) (list-fn2 place-ball EMPTY-SCENE lst))
Common List Function #2: foldr (start at right)

```scheme
;; foldr: (X Y -> Y) Y Listof<X> -> Y
(define (foldr fn initial lst)
  (cond
   [(empty? lst) initial]
   [else (fn (first lst) (foldr fn initial (rest lst)))]))

;; sum-lst: ListofInt -> Int
(define (sum-lst lst) (foldr + 0 lst))

;; render-world: ListofBall-> Image
(define (render-world lst) (foldr place-ball EMPTY-SCENE lst))
```

Function recurs and builds up fn calls until it gets to the end

Then they are evaluated, last one first
Common List Function #2: foldr

;; foldr: (X ... Y -> Y) Y Listof<X> ... -> Y

Racket version can also take multiple lists

(\(\text{foldr proc init lst ...+)} \rightarrow \text{any/c}\)

\begin{align*}
\text{proc : procedure?} \\
\text{init : any/c} \\
\text{lst : list?}
\end{align*}
Is it ok to always start at the right?

For some functions, order doesn’t matter, but for others, it does?

\[(\text{foldr} + 0 \ (\text{list} \ 1 \ 2 \ 3)) = (1 + (2 + (3 + 0)))\]

\[(1 + (2 + (3 + 0))) = ((1 + 0) + 2) + 3\]

\[(1 - (2 - (3 - 0))) = (((1 - 0) - 2) - 3) \neq \]
Need List Function #2b: \texttt{foldl} (start from left)

Challenge:

\begin{itemize}
\item Change \texttt{foldr} to \texttt{foldl}
\item so that the function \textbf{is applied from the left} (first element first)
\end{itemize}

\begin{verbatim}
(define (foldr fn initial lst)
  (cond
   [(empty? lst) initial]
   [else (fn (first lst) (foldr fn initial (rest lst)))]))

(define (foldl fn initial lst)
  (cond
   [(empty? lst) ....]
   [else .... (first lst) .... (foldl fn initial (rest lst)) ....]])
\end{verbatim}
Next time: Other common list functions

- Filter
- Find
- Reverse
- append
Check-In Quiz 10/11

on gradescope

(due 1 minute before midnight)