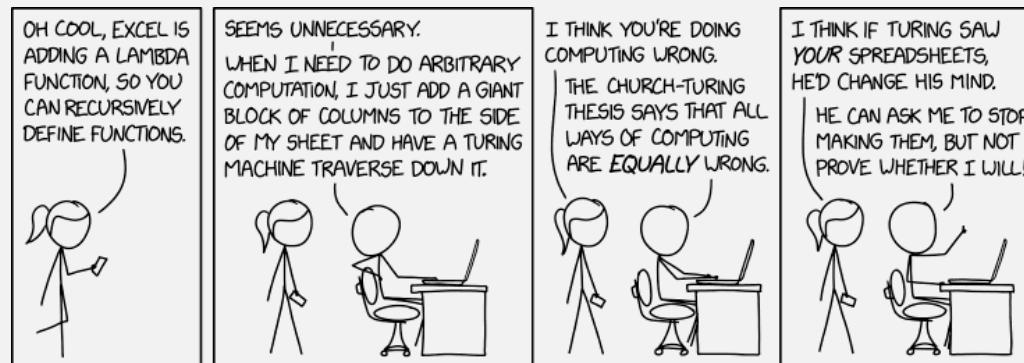


UMass Boston Computer Science
CS450 High Level Languages (section 2)

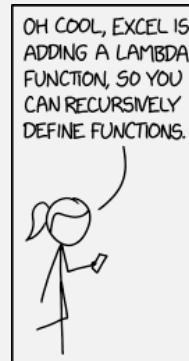
Function “Arithmetic” and the Lambda Calculus

Wednesday, October 18, 2023

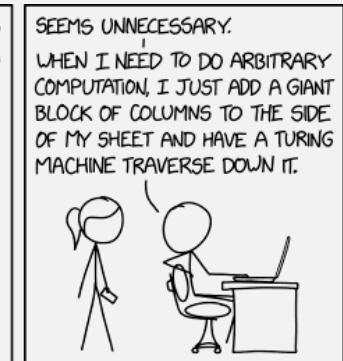


Logistics

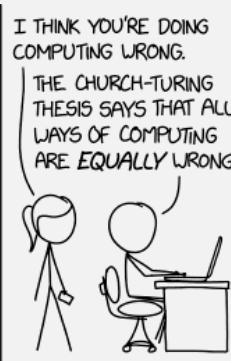
- HW 4 out
 - due: Sun 10/22 11:59 pm EST



OH COOL, EXCEL IS
ADDING A LAMBDA
FUNCTION, SO YOU
CAN RECURSIVELY
DEFINE FUNCTIONS.



SEEMS UNNECESSARY.
WHEN I NEED TO DO ARBITRARY
COMPUTATION, I JUST ADD A GIANT
BLOCK OF COLUMNS TO THE SIDE
OF MY SHEET AND HAVE A TURING
MACHINE TRAVERSE DOWN IT.



I THINK YOU'RE DOING
COMPUTING WRONG.
THE CHURCH-TURING
THESIS SAYS THAT ALL
WAYS OF COMPUTING
ARE EQUALLY WRONG.



I THINK IF TURING SAW
YOUR SPREADSHEETS,
HE'D CHANGE HIS MIND.
HE CAN ASK ME TO STOP
MAKING THEM, BUT NOT
PROVE WHETHER I WILL!

Common List Function #1: map

```
;; map: (X -> Y) Listof<X> -> Listof<Y>
;; Produces a list resulting from applying
;; a given fn to each element of a given lst
```

```
(define (map fn lst)
  (cond
    [(empty? lst) empty]
    [else (cons (fn (first lst))
                (map (rest lst))))]))
```

function “**application**”
(in high-level languages)
= function “**call**” (in
imperative languages)

`(map proc lst ...+) → list?`

procedure

Applies *proc* to the elements of the *lst*s from the first elements to the last. The *proc* argument must accept the same number of arguments as the number of supplied *lst*s, and all *lst*s must have the same number of elements. The result is a list containing each result of *proc* in order.

Examples:

```
> (map (lambda (number1 number2)
              (+ number1 number2))
        '(1 2 3 4)
        '(10 100 1000 10000))
      '(11 102 1003 10004)
```

RACKET’s map takes
multiple lists

map in other high-level languages

Array.prototype.map()

The `map()` method of `Array` instances creates a new array populated with the results of calling a provided function on every element in the calling array.

JavaScript Demo: Array.map()

```
1 const array1 = [1, 4, 9, 16];
2
3 // Pass a function to map
4 const map1 = array1.map(x => x * 2);
5
6 console.log(map1);
7 // Expected output: Array [2, 8, 18, 32]
```

Lambda
("arrow function expression")

Python3

```
# Add two lists using map and lambda

numbers1 = [1, 2, 3]
numbers2 = [4, 5, 6]

result = map(lambda x, y: x + y, numbers1, numbers2)
print(list(result))
```

lambda

Common List Function #2: foldl / foldr

```
;; foldr: (X Y -> Y) Y Listof<X> -> Y
;; Computes a single value from given list, determined by given fn and initial val.
;; fn is applied to each list element, last-element-first
(define (foldr fn initial lst)
  (cond
    [(empty? lst) initial]
    [else (fn (first lst) (foldr fn initial (rest lst))))])
```

$$\begin{array}{|c|} \hline (1 + (2 + (3 + 0))) \\ \hline (1 - (2 - (3 - 0))) \\ \hline \end{array}$$

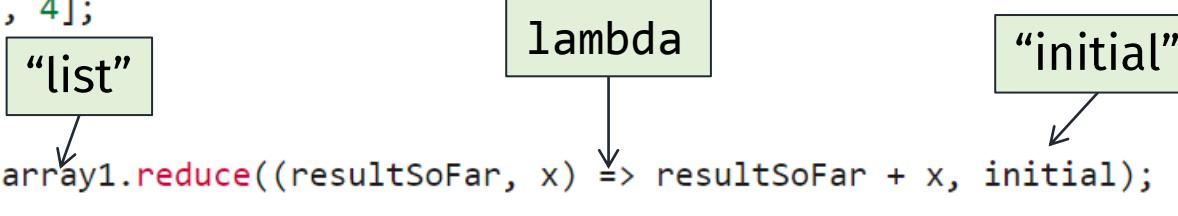
```
;; foldl: (X Y -> Y) Y Listof<X> -> Y
;; Computes a single value from given list, determined by given fn and initial val.
;; fn is applied to each list element, first-element-first
(define (foldl fn result-so-far lst)
  (cond
    [(empty? lst) result-so-far]
    [else (foldl fn (fn (first lst) result-so-far) (rest lst))]))
```

$$\begin{array}{|c|} \hline (((1 + 0) + 2) + 3) \\ \hline (((1 - 0) - 2) - 3) \\ \hline \end{array}$$

fold (reduce) in other high-level languages

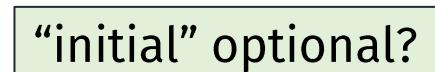
JavaScript Demo: Array.reduce()

```
1 const array1 = [1, 2, 3, 4];
2
3 // 0 + 1 + 2 + 3 + 4
4 const initialValue = 0;
5 const sumWithInitial = array1.reduce((resultSoFar, x) => resultSoFar + x, initialValue);
6
7 console.log(sumWithInitial);
8 // Expected output: 10
9
```



JavaScript Demo: Array.reduceRight()

```
1 const array1 = [
2   [0, 1],
3   [2, 3],
4   [4, 5],
5 ];
6
7 const result = array1.reduceRight((resultSoFar, x) => resultSoFar.concat(x));
8
9 console.log(result);
10 // Expected output: Array [4, 5, 2, 3, 0, 1]
11
```



Fold “dual”: build-list

(**build-list** *n proc*) → list?

n : exact-nonnegative-integer?

proc : (exact-nonnegative-integer? . -> . any)

procedure

Creates a list of *n* elements by applying *proc* to the integers from 0 to (**sub1** *n*) in order. If *lst* is the resulting list, then (**list-ref** *lst i*) is the value produced by (*proc i*).

Examples:

```
> (build-list 10 values)
'(0 1 2 3 4 5 6 7 8 9)
> (build-list 5 (lambda (x) (* x x)))
'(0 1 4 9 16)
```

```
(build-list 4 add1)
```

```
;; = (map add1 (list 0 1 2 3))
```

```
;; = (list 1 2 3 4)
```

Fold “alternative”: `apply` (and variable-arity fns)

```
(foldl + 0 (list 1 2 3 4)) ; = (+ (+ (+ (+ 1 0) 2) 3) 4)) = 10
```



```
(apply + (list 1 2 3 4)) ; = (+ 1 2 3 4) = 10
```

```
(apply string-append (list "a" "b" "cd")) ; = "abcd"
```

- `apply` applies its function arg to the contents of its list arg
- function arg to `apply` must accept:
of arguments = length of list arg

Common list function #3: filter

```
;; filter: Listof<X> (X -> Boolean) -> Listof<X>
;; Returns a list containing elements of given list
;; for which the given predicate returns true
```

```
(define (filter lst pred?)
  (cond
    [(empty? lst) empty]
    [else (if (pred? (first lst))
               (cons (first lst) (filter (rest lst)))
               (filter (rest lst))))]))
```

filter in other high-level languages

JavaScript Demo: Array.filter()

```
1 const words = ['spray', 'limit', 'elite', 'exuberant', 'destruction', 'present'];
2
3 const result = words.filter((word) => word.length > 6);
4
5 console.log(result);
6 // Expected output: Array ["exuberant", "destruction", "present"]
7
```

Common list function #3: filter

```
;; filter: Listof<X> (X -> Boolean) -> Listof<X>
;; Returns a list containing elements of given list for which the given predicate returns true
```

```
(define (filter lst pred?)
  (cond
    [(empty? lst) empty]
    [else (if (pred? (first lst))
               (cons (first lst) (filter (rest lst) pred?))
               (filter (rest lst))))]))
```

lambda rules:

- Can skip the **design recipe** steps, BUT
- **name, description, and signature** must be "obvious"
- **code is arithmetic only**
- otherwise, create standalone function **define**

```
;; smaller-than: Listof<Int> Int -> Listof<Int>
;; Returns a list containing elements of given list less than the given int
```

```
(define (smaller-than lst thresh)
  (filter (lambda (x) (< x thresh)) lst))
```

lambda creates an anonymous “inline” function (expression)

Functions as Values

- In high-level languages, functions are no different from other values (e.g., numbers)
- They can be passed around, or be the result of a function

```
;; make< : Int -> (Int -> Bool)
;; makes a function that returns true
;; for values less than the given thresh value
```

```
(define (make< thresh)
  (lambda (x) (< x thresh)))
```

```
(define (smaller-than lst thresh)
  (filter (make< thresh)) lst)
```

- lambda is just one way to “make” functions
- We can also do “arithmetic” with functions

Currying

- A “curried” function is “partially” applied to some (but not all) args
- Result is another function

```
(curry < 4)
;; = a function that returns true when given a number less than 4
```

```
(define (smaller-than lst thresh)
  (filter (lambda (x) (< x thresh)) lst))
```



```
(define (smaller-than lst thresh)
  (filter (curry > thresh)) lst))
```

History Lesson: Haskell B. Curry



- Mathematician / Logician
- Born in Millis, MA, in year 1900
- “currying” functions is named after him
- and also, the “Haskell” (functional) programming language
- Invented “combinatory logic”,
i.e., a system of function “arithmetic”

Currying

```
(define (smaller-than lst thresh)
  (filter (lambda (x) (< x thresh)) lst))
```

```
(define (smaller-than lst thresh)
  (filter (curry > thresh)) lst))
```

```
(define (smaller-than lst thresh)
  (filter (curryr < thresh)) lst))
```

NOTE: First argument gets curried first

Composing Functions

- compose combines multiple functions into one function
 - last one is applied first

```
(compose sqrt add1)
;; = a function that first applies add1 to its argument, then sqrt
```

```
((compose sqrt add1) 8) ; = 3
```

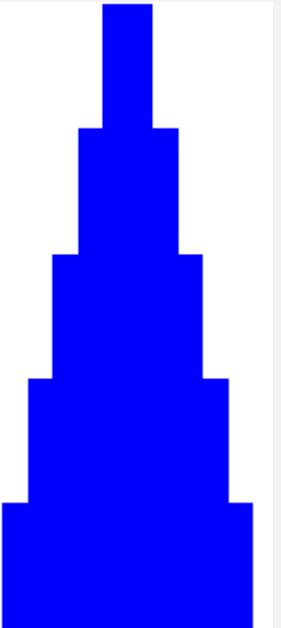
Composing Functions

- compose combines multiple functions into one function
 - last one is applied first

```
6 apply  
5 above  
1 (build-list ; = (list 0 1 2 3 4)  
5  
4 (compose  
3 (curryr square "solid" "blue")  
2 (curry * 20) ; = (list 20 40 60 80 100))  
add1))) ; = (list 1 2 3 4 5)
```

```
; = (above (square 20 "solid" "blue")  
          (square 40 "solid" "blue")  
          (square 60 "solid" "blue")  
          (square 80 "solid" "blue")  
          (square 100 "solid" "blue"))
```

```
; = (list (square 20 "solid" "blue")  
          (square 40 "solid" "blue")  
          (square 60 "solid" "blue")  
          (square 80 "solid" "blue")  
          (square 100 "solid" "blue"))
```



The Lambda (λ) Calculus

- A “programming language” consisting of only:
 - Lambda
 - Function application
- Equivalent in “computational power” to
 - Turing Machines
 - Your favorite programming language!

History Lesson: Alonzo Church

- Mathematician, logician, computer scientist
- Invented the lambda calculus
- And (half of) Church-Turing Thesis
 - Any function that can be “computed” has an equivalent Turing Machine
 - And an equivalent program in the lambda calculus
 - so, a Turing Machine = a lambda



Church Numerals

```
;; A ChurchNum is a function with two arguments:  
;; "fn" : a function to apply  
;; "base" : a base ("zero") value to apply to  
;;  
;; For a specific number, its "Church" representation  
;; applies the given function that number of times
```

```
(define czero  
  (lambda (f base) base))
```

Function applied zero times

```
(define cone  
  (lambda (f base) (f base)))
```

Function applied one times

```
(define ctwo  
  (lambda (f base) (f (f base))))
```

Function applied two times

```
(define cthree  
  (lambda (f base) (f (f (f base))))))
```

Function applied three times

Church “Add1”

```
;; cplus1 : ChurchNum -> ChurchNum  
;; “Adds” 1 to the given Church num
```

```
(define cplus1  
  (lambda (n)  
    (lambda (f base)  
      (f (n f base))))))
```

Input ChurchNum

Returns a ChurchNum ...

(we know “n” will apply f n times)

... that adds an extra application of f

```
(define czero  
  (lambda (f base) base))
```

```
(define cone  
  (lambda (f base) (f base)))
```

```
(define ctwo  
  (lambda (f base) (f (f base))))
```

```
(define cthree  
  (lambda (f base) (f (f (f base))))))
```

Church Addition

```
;; cplus : ChurchNum ChurchNum -> ChurchNum  
;; “Adds” the given ChurchNums together
```

```
(define cplus  
  (lambda (m n)  
    (lambda (f base)  
      (m f (n f base)))))
```

Input ChurchNums

Returns a ChurchNum ...

(we know “n” will apply f n times)

... that adds “m” extra application of f

```
(define czero  
  (lambda (f base) base))
```

```
(define cone  
  (lambda (f base) (f base)))
```

```
(define ctwo  
  (lambda (f base) (f (f base))))
```

```
(define cthree  
  (lambda (f base) (f (f (f base))))))
```

Code Demo 1

Church Booleans

```
;; A ChurchBool is a function with two arguments,  
;; where the representation of:  
;; “true” returns the first arg, and  
;; “false” returns the second arg
```

```
(define ctrue  
  (lambda (a b) a))
```

Returns first arg

```
(define cfalse  
  (lambda (a b) b))
```

Returns second arg

Review: “And”

The truth table of $A \wedge B$:

A	B	$A \wedge B$
True	True	True
True	False	False
False	True	False
False	False	False

When $A = \text{True}$,
then $\text{And}(A, B) = B$

When $A = \text{False}$,
then $\text{And}(A, B) = A$

Church “And”

The truth table of $A \wedge B$:

A	B	$A \wedge B$
True	True	True
True	False	False
False	True	False
False	False	False

When $A = \text{True}$,
want $\text{And}(A, B) = B$

When $A = \text{False}$,
want $\text{And}(A, B) = A$

`;; cand: ChurchBool ChurchBool-> ChurchBool
;; “ands” the given ChurchBools together`

```
(define cand
  (lambda (A B)
    (A B A)))
```

```
(define ctrue
  (lambda (a b) a))
```

;; if $A = \text{ctrue}$

;; then $(A B A) = B$

;; want $(\text{cand } A B) = B$

(Returns first arg)

```
(define cfalse
  (lambda (a b) b))
```

;; if $A = \text{cfalse}$

;; then $(A B A) = A$

;; want $(\text{cand } A B) = A$

(Returns second arg)

Church “Or”

`;; cor: ChurchBool ChurchBool-> ChurchBool
;; “or” the given ChurchBools together`

A	B	$A \vee B$
True	True	True
True	False	True
False	True	True
False	False	False

When $A = \text{True}$, want $\text{Or}(A, B) = A$

When $A = \text{False}$, want $\text{Or}(A, B) = B$

```
(define cor
  (lambda (A B)
    (A A B)))
```

```
(define ctrue
  (lambda (a b) a))
```

```
;; if A = ctrue
;; then (A A B) = A   
;; want (cor A B) = A
```

(Returns first arg)

```
(define cfalse
  (lambda (a b) b))
```

```
;; if A = cfalso
;; then (A A B) = B   
;; want (cor A B) = B
```

(Returns second arg)

Code Demo 2

Church Pairs (Lists)

```
;; A ChurchPair<X,Y> 1-arg function, where  
;; the arg fn is applied to (i.e., "selects") the X and Y data values
```

```
;; ccons: X Y -> ChurchPair<X,Y>
```

```
(define ccons  
  (lambda (x y)  
    (lambda (get)  
      (get x y))))
```

```
(define cfirst  
  (lambda (cc)  
    (cc (lambda (x y) x))))
```

```
(define csecond  
  (lambda (cc)  
    (cc (lambda (x y) y))))
```

“Gets” the first item

“Gets” the second item

Code Demo 3

The Lambda Calculus

- A “programming language” consisting of only:
 - Lambda
 - Function application
- “Language” has:
 - Numbers
 - Booleans and conditionals
 - Lists
 - ...
 - Recursion?

Check-In Quiz 10/18

on gradescope

(due 1 minute before midnight)