Introduction to Compiler Construction in a Java World

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Errata

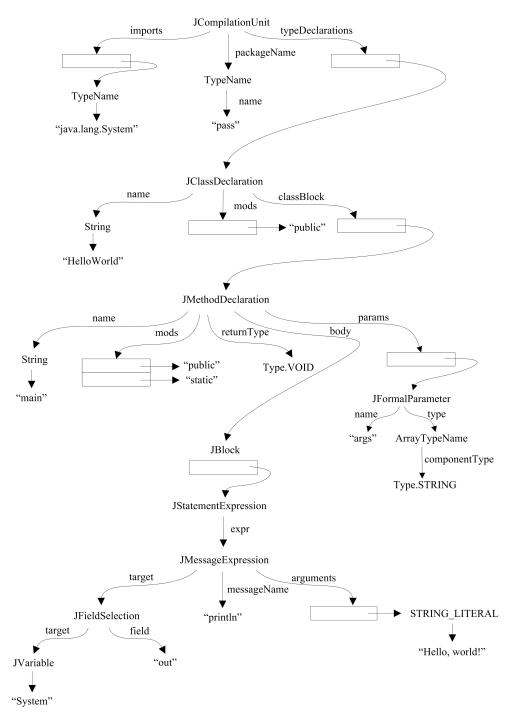
Here you can find a listing of known errors in our text. If you find others, please let us know about them at j--@cs.umb.edu. We appreciate your feedback.

Chapter 1: Compilation

On page 6, Figure 1.6 the "JVM Code" on the left of the figure must be "Source Language Program". Here is the updated figure.



Reported by Antoine Marchal on Jun 17, 2014



On page 14, Figure 1.9 is missing the formal parameters of the main() method. Here is the updated figure.

Reported by Bill Campbell on Jan 28, 2013

On page 18, the following line

For example, the j-- program ...

should be

For example, the j-- program . . .

Reported by Pierre Schaus on Feb 7, 2013

On page 19, the following code snippet

```
public class DivisionTest extends TestCase {
    ...
```

should be

```
package junit;
import junit.framework.TestCase;
import pass.Division;
public class DivisionTest extends TestCase {
    ...
```

Reported by Daisuke Tanaka on Feb 1, 2013

Chapter 2: Lexical Analysis

On page 41, the set of moves M currently given by

$$M = \{m(0, a) = 1, m(0, b) = 1, m(1, a) = 1, m(1, b) = 1, m(1, \epsilon) = 0, m(1, b) = 2\}$$

should be

$$M = \{m(0, a) = 1, m(1, a) = 1, m(1, b) = 1, m(1, \epsilon) = 0, m(1, b) = 2\}$$

Reported by Pierre Schaus on Feb 11, 2013

On page 46, the following line in Definition 2.6

 \ldots set of states S includes s and \ldots

should be

 \ldots set of states S includes S and \ldots

Reported by Pierre Schaus on Feb 7, 2013

Chapter 3: Parsing

On page 59, the following code snippet

package pass;f
 ...

should be

```
package pass;
...
```

Reported by Pierre Schaus on Feb 14, 2013

On page 88, the following line in Algorithm 3.6

where $X_j ::= \beta_1 | \beta_2 | \dots | \beta_k$ are the current rules defining X_i

should be

where $X_j ::= \beta_1 |\beta_2| \dots |\beta_k$ are the current rules defining X_j

Reported by Pierre Schaus on Feb 14, 2013

Chapter 5: JVM Code Generation

The tables on page 187 and 189 are incorrect. Here are the updated tables.

	x	a[i]	o.f	C.sf
lhs = y	iload y'	aload a'	aload o'	iload y'
	[dup]	iload i'	iload y	[dup]
	istore x'	iload y'	[dup_x1]	putstatic sf
		[dup_x2]	putfield f	
		iastore		
lhs += y	iload x'	aload a'	aload o'	getstatic sf
	iload y'	iload i'	dup	iload y'
	iadd	dup2	getfield f	iadd
	[dup]	iaload	iload y'	[dup]
	istore x'	iload y'	iadd	putstatic sf
		iadd	[dup_x1]	
		[dup_x2]	putfield f	
		iastore		
++lhs	iinc x',1	aload a'	aload o'	getstatic sf
	[iload x']	iload i'	dup	iconst_1
		dup2	getfield f	iadd
		iaload	iconst_1	[dup]
		iconst_1	iadd	putstatic sf
		iadd	[dup_x1]	
		[dup_x2]	putfield f	
		iastore		
lhs	[iload x']	aload a'	aload o'	getstatic sf
	iinc x',-1	iload i'	dup	[dup]
		dup2	getfield f	iconst_1
		iaload	[dup_x1]	isub
		[dup_x2]	iconst_1	putstatic sf
		iconst_1	isub	
		isub	putfield f	
		iastore		

	x	a[i]	o.f	C.sf
codegenLoadLhsLvalue()	[none]	aload a'	aload o'	[none]
		iload i'		
codegenLoadLhsRvalue()	iload x'	dup2	dup	getstatic sf
		iaload	getfield f	
<pre>codegenDuplicateRvalue()</pre>	dup	dup_x2	dup_x1	dup
codegenStore()	istore x'	iastore	putfield f	putstatic sf

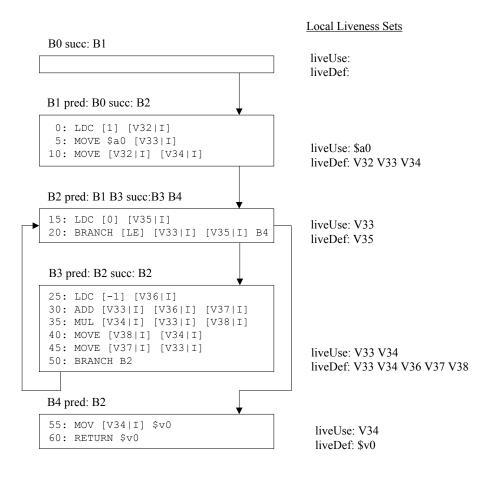
Reported by Bill Campbell on April 15, 2013

Chapter 7: Register Allocation

In Figure 7.1 on page 247 and Figure 7.3 on page 250, the link from block B2 to block B3 is missing. Here are the updated figures.

B0 succ: B1

0:	red: B0 succ: B2
B2 p	red: B1 B3 succ:B3 B4
	LDC [0] [V35 I] BRANCH [LE] [V33 I] [V35 I] B4
В3 р	ored: B2 succ: B2
25: 30: 35: 40: 45:	Dred: B2 succ: B2
25: 30: 35: 40: 45: 50:	LDC [-1] [V36 I] ADD [V33 I] [V36 I] [V37 I] MUL [V34 I] [V33 I] [V38 I] MOVE [V38 I] [V34 I] MOVE [V37 I] [V33 I]



Reported by Josef Joller on Dec 15, 2013

Algorithm 7.9 uses ! for logical not, which is confusing since ! also represents logical not in j--. Here is the corrected version, with ! replaced by **not**.

Algorithm 1 Graph Coloring Register Allocation

Input: The control-flow graph g for a method with LIR that makes use of virtual registers **Output:** The same g but with virtual registers replaced by physical registers $registersAssignedSuccessfully \leftarrow \texttt{false}$ repeat repeat buildIntervals() buildInterferenceGraph() until not coalesceRegistersSuccessful() buildAdjacencyLists() computeSpillCosts() pruneGraph() $registersAssignedSuccessfully \leftarrow assignRegisters()$ if not *registersAssignedSuccessfully* then generateSpillCode() end if until registersAssignedSuccessfully

Reported by Bill Campbell on Nov 30, 2014