Consistently better than SGD

Best regret bound

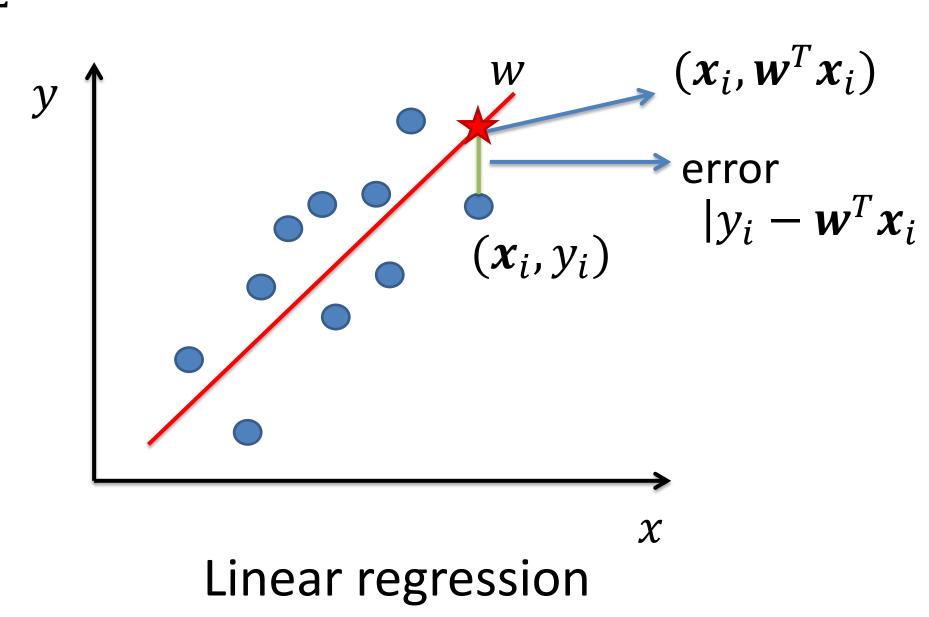
Ready to use

Code: http://www.cs.umb.edu/~yangmu/code/CSGD.zip

1. The Least Squares problem

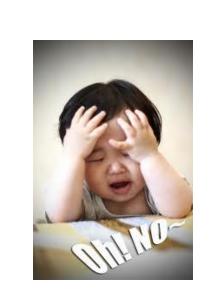
Objective function:

 $\min_{\pmb{w}} \sum_{i=1}^{1} \|y_i - \pmb{w}^T \pmb{x}_i\|_2^2$, where $\pmb{x}_i \in \mathbb{R}^d$, $y_i \in \mathbb{R}^1$ and $\pmb{w} \in \mathbb{R}^d$



2. When Least Squares problem meets Large Scale

Use closed form solution: $\mathbf{w} = (\mathbf{X}^T \mathbf{X})^{-1} \mathbf{X}^T \mathbf{y}$?



It's $O(nd^2)$!

Note: Any algorithm with time complexity greater than O(nd)is not applicable in large scale high dimension cases.

Stochastic Gradient Descent (SGD) with O(d) time each iteration is an appealing approach, which takes the form

$$\mathbf{w}_{t+1}^* = \arg\min_{\mathbf{w}} \frac{1}{t} \sum_{i=1}^t \frac{1}{2} ||y_i - \mathbf{w}^T \mathbf{x}_i||_2^2,$$

where $l(\mathbf{w}, \mathbf{x}_i, y_i) = \frac{1}{2} ||y_i - \mathbf{w}^T \mathbf{x}_i||_2^2$ is the loss function at step i, denoted as l(w) for short.

SGD update rule is:

$$\boldsymbol{w}_{t+1} = \boldsymbol{w}_t - \eta_t \boldsymbol{g}_t,$$

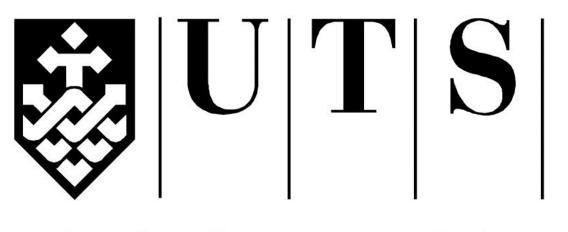
where $\boldsymbol{g}_t = \partial l(\boldsymbol{w})$.

Constrained Stochastic Gradient Descent for Large-scale Least Squares Problem

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 $w^*^T \overline{x} = \overline{y}$

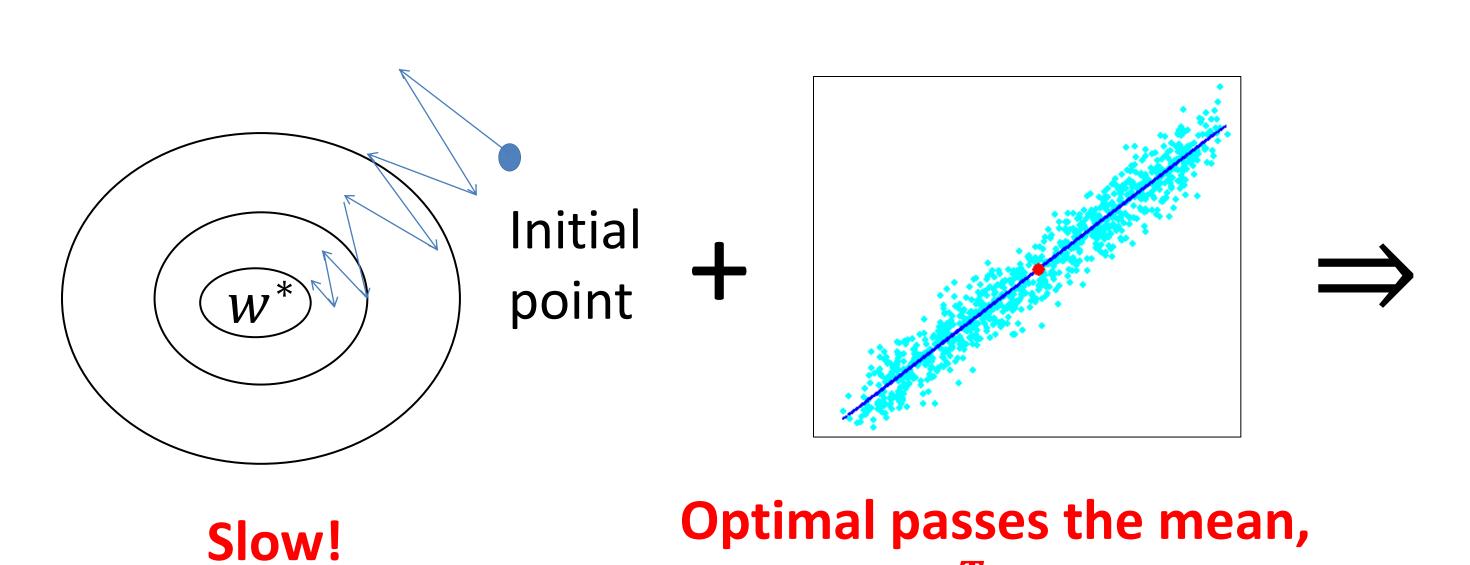


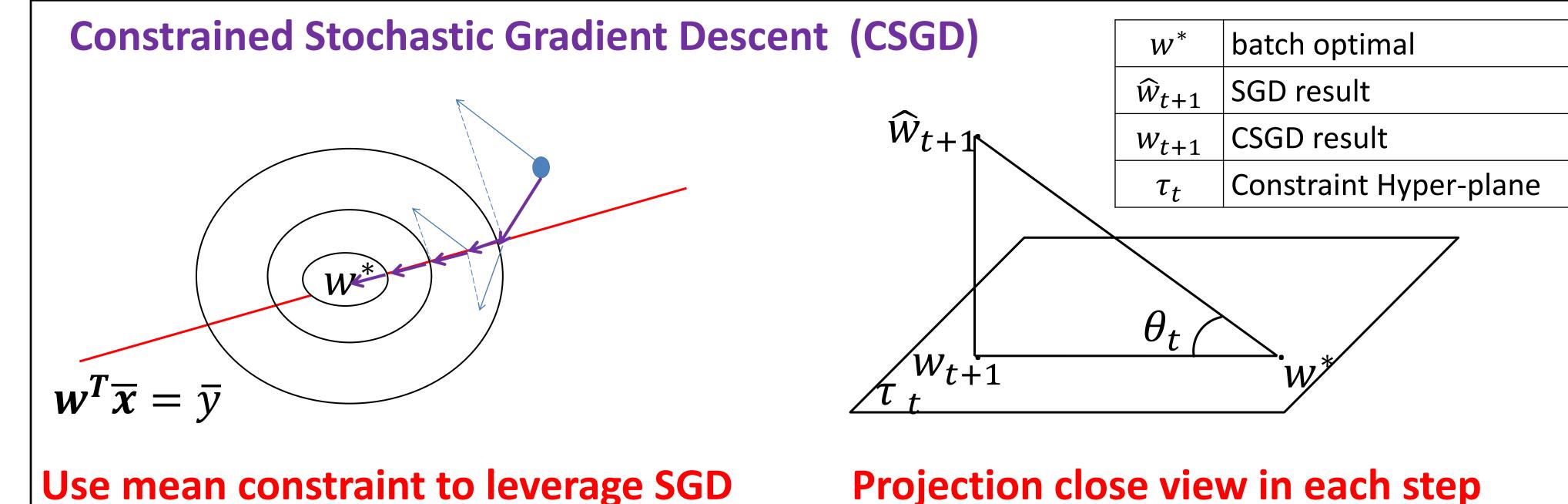


University of Technology, Sydney

3. Motivation

Zig-Zag property of SGD The beauty of Least Squares





Remarks:

1. performs consistently better than SGD in terms of batch optimal. 2. has $O(\log T)$ regret bound $R_G(T) \le \frac{G^2}{2H}(1 + \log T)$ 3. extensible.

4. Algorithm

CSGD:

 $w_{t+1}^* = \operatorname{arg\,min}_{w} \frac{1}{t} \sum_{i=1}^{t} \frac{1}{2} \|y_i - w^T x_i\|_2^2$, s,t. $w^T \overline{x}_t = \overline{y}_t$. The update rule $w_{t+1} = P_t(w_t - \eta_t g_t) + r_t$, where $P_t = I - \frac{\overline{x}_t \overline{x}_t}{\|\overline{x}_t\|_2^2}$, $r_t = \frac{\overline{y}_t}{\|\overline{x}_t\|_2^2} \overline{x}_t$

WAIT! Is $P_t \in \mathbb{R}^{d \times d}$? **YES**, but it is rank ONE!

Therefore, we still have a O(d) time complexity algorithm each iteration: $w_{t+1} = w_t - \eta_t g_t - \overline{x}_t (\overline{x}_t^T (w_t - \eta_t g_t)) / \|\overline{x}_t\|_2^2 + r_t$.

5. Experiments

