Finite Automata and Regular Languages (part II)

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UMB
1 Nondeterministic Automata
Definition

A nondeterministic finite automaton (ndfa) is a quintuple \( \mathcal{M} = (A, Q, \delta, q_0, F) \), where \( A \) is the input alphabet of \( \mathcal{M} \), \( Q \) is a finite set of states, \( \delta : Q \times A \rightarrow \mathcal{P}(Q) \) is the transition function, \( q_0 \in Q \) is the initial state, and \( F \subseteq Q \) is the set of final states of \( \mathcal{M} \). We assume \( A \cap Q = \emptyset \).
Example
Consider the ndfa

$$M = (\{a, b\}, \{q_0, q_1, q_2, q_3, q_4\}, \delta, q_0, \{q_1, q_3\})$$

whose transition function is defined by the table:

<table>
<thead>
<tr>
<th>Input</th>
<th>State</th>
<th>q₀</th>
<th>q₁</th>
<th>q₂</th>
<th>q₃</th>
<th>q₄</th>
</tr>
</thead>
<tbody>
<tr>
<td>a</td>
<td>{q₁, q₂}</td>
<td>∅</td>
<td>{q₃}</td>
<td>∅</td>
<td>∅</td>
<td>∅</td>
</tr>
<tr>
<td>b</td>
<td>{q₀}</td>
<td>{q₃}</td>
<td>{q₄}</td>
<td>∅</td>
<td>∅</td>
<td></td>
</tr>
</tbody>
</table>

Note the presence of pairs \((q, a)\) such that \(\delta(q, a) = \emptyset\). We refer to such pairs as blocking situations of \(M\). For instance, \((q_1, a)\) is a blocking situation of \(M\).
Extending the transition function for an ndfa

As in the case of the dfa, we can extend the ndfa’s transition function $\delta$, defined on single characters, to $\delta^*$, defined on words. Starting from the transition function $\delta$, we define the function $\delta^*: Q \times A^* \rightarrow \mathcal{P}(Q)$ as follows:

\[
\delta^*(q, \lambda) = \{q\}
\]

\[
\delta^*(q, xa) = \bigcup_{q' \in \delta^*(q,x)} \delta(q', a)
\]
If $\mathcal{M} = (A, Q, \delta, q_0, F)$ is an ndfa, then the graph of $\mathcal{M}$ is the marked, directed multigraph $\mathcal{G}(\mathcal{M})$, whose set of vertices is the set of states $Q$.

The set of edges of $\mathcal{G}(\mathcal{M})$ consists of all pairs $(q, q')$ such that $q' \in \delta(q, a)$ for some $a \in A$; an edge $(q, q')$ is labeled by the input symbol $a$, where $q' \in \delta(q, a)$.

The initial state $q_0$ is denoted by an incoming arrow with no source, and the final states are circled.

If $q' \in \delta^*(q, \alpha)$, then there is a path in the graph $\mathcal{G}(\mathcal{M})$ labeled by $\alpha$ that leads from $q$ to $q'$.
Nondeterministic Automata

Diagram of a nondeterministic automaton with states $q_0$, $q_1$, $q_2$, $q_3$, and $q_4$, and transitions labeled with $a$ and $b$. The diagram shows self-loops on $q_0$ and $q_2$, and transitions from $q_0$ to $q_1$ and $q_2$, from $q_1$ to $q_3$, from $q_2$ to $q_4$, and from $q_3$ to $q_4$.
Comparing dfas and ndfas

- In the graph of a dfa \( M = (A, Q, \delta, q_0, F) \) you must have exactly one edge emerging from every state \( q \) and for every input symbol \( a \in A \).

- In the graph of an ndfa \( M = (A, Q, \delta, q_0, F) \) you may have states where no arrow emerges, or states where several arrow labeled with the same symbol emerge. Also, not every symbol needs to appear as a label of an emerging edge.
Definition

The language accepted by the ndfa $\mathcal{M} = (A, Q, \delta, q_0, F)$ is

$$L(\mathcal{M}) = \{x \in A^* \mid \delta^*(q_0, x) \cap F \neq \emptyset\}.$$ 

In other words, $x \in L(\mathcal{M})$ if there exists a path in the graph of $\mathcal{M}$ labeled by $x$ that leads from the initial state into one of the final states. Note that it is not necessary that all paths labeled by $x$ lead to a final state; the existence of one such path suffices to put $x$ into the language $L(\mathcal{M})$. 
Note that $ab \in L(M)$ because of the existence of the path $(q_0, q_1, q_3)$ labeled by this word and the fact that $q_3$ is a final state. On the other hand, $(q_0, q_2, q_4)$ is another path labeled by $x$ but $q_4 \notin F$. 
Example (cont’d)

This ndfa is simple enough to allow an easy identification of all types of words in $L(M)$:

1. The final state $q_1$ can be reached by applying an arbitrary number of $bs$ followed by an $a$.

2. The final state $q_3$ can be reached by a path of the form $(q_0, \ldots, q_0, q_1, q_3)$, that is by a word of the form $b^k ab$ for $k \in \mathbb{N}$.

3. The same final state $q_3$ can be reached via $q_2$. Input words that allow this transition have the form $b^k aa$ for $k \in \mathbb{N}$.

Thus, we have

$$L(M) = \{ b \}^* a \cup \{ b \}^* ab \cup \{ b \}^* aa = \{ b \}^* \{ a, ab, aa \}.$$
Example

Consider an alphabet $A = \{a_0, \ldots, a_{n-1}\}$ and a binary relation $\rho \subseteq A \times A$. The language

$$L_\rho = \{a_{i_0} \cdots a_{i_p} \mid p \in \mathbb{N}, (a_{i_j}, a_{i_{j+1}}) \in \rho \text{ for } 0 \leq j \leq p - 1\}$$

is accepted by the ndfa $M_\rho = (A, Q, \delta, q_0, F)$, where $Q = \{q, q_0, \ldots, q_{n-1}\}$, $F = \{q_0, \ldots, q_{n-1}\}$, and $\delta$ is given by

- $\delta(q, a_i) = \{q_i\}$ for $0 \leq i \leq n - 1$;
- for every $i, j$ such that $0 \leq i, j \leq n - 1$,

$$\delta(q_i, a_j) = \{q_j \in Q \mid (a_i, a_j) \in \rho\}.$$
Example (cont’d)

Note that if \((a_i, a_j) \notin \rho\), then \((q_i, a_j)\) is a blocking situation. The existence of these blocking situations is precisely what makes this device a nondeterministic automaton.
We show that $L_{\rho} = L(M_{\rho})$. Let $x = a_{i_0} \cdots a_{i_p}$ be a word in $L_{\rho}$ with $p \geq 0$. We prove by induction on $p = |x| - 1$ that $x \in L(M_{\rho})$ and that $\delta^*(q, x) = \{q_{i_p}\}$. The base case, $p = 0$, is immediate, since the condition $(a_{i_j}, a_{i_{j+1}}) \in \rho$ for $0 \leq j \leq p - 1$ is vacuous.

Suppose that the statement holds for words in $L_{\rho}$ of length at most $p$ and let $x = a_{i_0} \cdots a_{i_p}$ be a word in $L_{\rho}$ of length $p + 1$. It is clear that the word $y = a_{i_0} \cdots a_{i_{p-1}}$ belongs to $L_{\rho}$. By the inductive hypothesis, $y \in L(M_{\rho})$ and $\delta^*(q, y) = \{q_{i_{p-1}}\}$. Since $(a_{i_{p-1}}, a_{i_p}) \in \rho$ (by the definition of $L_{\rho}$), we have $\delta(q_{i_{p-1}}, a_{i_p}) = \{q_{i_p}\}$, so

$$q_{i_p} \in \bigcup_{q' \in \delta(q, y)} \delta(q', a_{i_p}) = \delta^*(q, ya_{i_p}) = \delta^*(q, x).$$

Therefore, $x \in L(M_{\rho})$. 
To prove the converse inclusion $L(M_\rho) \subseteq L_\rho$ we use an argument by induction on $|x| \geq 1$, where $x$ is a word from $L(M_\rho)$, to show that if $x = a_{i_0} \cdots a_{i_p} \in L(M_\rho)$, then $\delta^*(q, x) = \{q_{i_p}\}$ and $x \in L_\rho$. Again, the base case is immediate.

Suppose that the statement holds for words in $L(M_\rho)$ of length less than $p + 1$ that belong to $L_\rho$, and let $x = a_{i_0} \cdots a_{i_p}$ be a word in $L(M_\rho)$ of length $p + 1$. If $y = a_{i_0} \cdots a_{i_p-1}$, it is easy to see that $y \in L(M_\rho)$ because no blocking situation may arise in $M_\rho$ while the symbols of $y$ are read. Therefore, by the inductive hypothesis, $y \in L_\rho$ and $\delta^*(q, y) = \{q_{i_p-1}\}$.

Further, since $\delta(q_{i_p-1}, a_{i_p}) \neq \emptyset$, it follows that $(a_{i_p-1}, a_{i_p}) \in \rho$, so $x \in L_\rho$, and $\delta^*(q, x) = q_{i_p}$.

Thus, $L_\rho$ is accepted by the ndfa $M_\rho$. 
Let $\mathcal{M} = (A, Q, \delta, q_0, F)$ be a nondeterministic automaton, and let $\Delta : \mathcal{P}(Q) \times A \rightarrow \mathcal{P}(Q)$ be defined by

$$\Delta(S, a) = \bigcup_{q \in S} \delta(q, a)$$  \hspace{1cm} (1)

for every $S \subseteq Q$ and $a \in A$. Starting from $\Delta$, we define $\Delta^* : \mathcal{P}(Q) \times A^* \rightarrow \mathcal{P}(Q)$ in the manner used for the transition functions of deterministic automata. Namely, we define

$$\Delta^*(S, \lambda) = S$$  \hspace{1cm} (2)

$$\Delta^*(S, xa) = \Delta(\Delta^*(S, x), a)$$  \hspace{1cm} (3)

for every $S \subseteq Q$ and $a \in A$. 
Lemma

The functions $\Delta$ and $\Delta^*$ defined above satisfy the following properties:

1. For every family of sets $\{S_0, \ldots, S_{n-1}\}$ and every $a \in A$, we have:

   $$\Delta \left( \bigcup_{0 \leq i \leq n-1} S_i, a \right) = \bigcup_{0 \leq i \leq n-1} \Delta(S_i, a).$$

2. For every set $S \subseteq Q$ and $x \in A^*$ we have

   $$\Delta^*(S, x) = \bigcup_{q \in S} \delta^*(q, x).$$
Proof

The first part of the lemma is immediate, because

\[ \Delta\left(\bigcup_{0 \leq i \leq n-1} S_i, a\right) = \bigcup\{\delta(q, a) \mid q \in \bigcup_{0 \leq i \leq n-1} S_i\} \]
\[ = \bigcup_{0 \leq i \leq n-1}\{\delta(q, a) \mid q \in S_i\} \]
\[ = \bigcup_{0 \leq i \leq n-1} \Delta(S_i, a). \]
The argument for the second part of the lemma is by induction on $|x|$. For the basis step, we have $|x| = 0$, so $x = \lambda$, and $\Delta^*(S, \lambda) = S$, $\bigcup_{q \in S} \delta^*(q, \lambda) = \bigcup_{q \in S} \{q\} = S$.

Suppose that the argument holds for words of length $n$, and let $x = za$ be a word of length $n + 1$. We have

\[
\Delta^*(S, x) = \Delta^*(S, za) \\
= \Delta(\Delta^*(S, z), a) \\
= \Delta(\bigcup_{q \in S} \delta^*(q, z), a) \text{(by ind. hyp.)} \\
= \bigcup_{q \in S} \Delta(\delta^*(q, z), a) \\
= \bigcup_{q \in S} \bigcup_{r \in \delta^*(q, z)} \delta(r, a) = \bigcup_{q \in S} \delta^*(q, za) = \bigcup_{q \in S} \delta^*(q, x).
\]
Nondeterministic automata can be regarded as generalizations of deterministic automata in the following sense. If $\mathcal{M} = (A, Q, \delta, q_0, F)$ is a deterministic automaton, consider a nondeterministic automaton $\mathcal{M}' = (A, Q, \delta', q_0, F)$, where $\delta'(q, a) = \{\delta(q, a)\}$. It is easy to verify that for every $q \in Q$ and $x \in A^*$ we have $\delta'^*(q, x) = \{\delta^*(q, x)\}$. Therefore,

$$L(\mathcal{M}') = \{x \in A^* \mid \delta'^*(q_0, x) \cap F \neq \emptyset\}$$

$$= \{x \in A^* \mid \{\delta^*(q_0, x)\} \cap F \neq \emptyset\}$$

$$= \{x \in A^* \mid \delta^*(q_0, x) \in F\}$$

$$= L(\mathcal{M}).$$

In other words, for every deterministic finite automaton there exists a nondeterministic one that recognizes the same language.
Theorem

For every nondeterministic automaton, there exists a deterministic automaton that accepts the same language.
Proof

Let \( \mathcal{M} = (A, Q, \delta, q_0, F) \) be a nondeterministic automaton. Define the function \( \Delta \) as in the equality

\[
\Delta(S, a) = \bigcup_{q \in S} \delta(q, a),
\]

and consider the deterministic automaton

\[\mathcal{M}' = (A, P(Q), \Delta, \{q_0\}, \{S \mid S \subseteq Q \text{ and } S \cap F \neq \emptyset\}).\]

For every \( x \in A^* \) we have the following equivalent statements:

1. \( x \in L(M) \);
2. \( \delta^*(q_0, x) \cap F \neq \emptyset \);
3. \( \Delta^*(\{q_0\}, x) \cap F \neq \emptyset \);
4. \( x \in L(M') \).

This proves that \( L(M) = L(M') \).
Example
Consider the nondeterministic finite automaton

\[ M = (\{a, b\}, \{q_0, q_1, q_2\}, \delta, q_0, \{q_2\}) \]

whose graph is given below.
It is easy to see that the language accepted by this automaton is $A^* bbA^*$, that is the language that consists of all words that contain two consecutive $b$ symbols.

The graph of the nondeterministic automaton is simpler than the graph of the previous deterministic automaton; this simplification is made possible by the nondeterminism.
Graph of the Equivalent ndfa